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## Construction of self-reciprocal normal polynomials over finite fields of even characteristic

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**Abstract:** In this paper, a computationally simple and explicit construction of some sequences of normal polynomials and self-reciprocal normal polynomials over finite fields of even characteristic are presented.

**Key words:** Finite fields, normal polynomial, self-reciprocal

### 1. Introduction

Let  $\mathbb{F}_q$ , be the Galois field of order  $q = p^s$ , where  $p$  is a prime and  $s$  is a natural number, and  $\mathbb{F}_q^*$  be its multiplicative group. Let  $P(x)$  be a monic irreducible polynomial of degree  $n$  over  $\mathbb{F}_q$  and  $\beta$  be a root of  $P(x)$ . The field  $\mathbb{F}_q(\beta) = \mathbb{F}_{q^n}$  is an  $n$ -dimensional extension of  $\mathbb{F}_q$  and can be considered as a vector space of dimension  $n$  over  $\mathbb{F}_q$ . The Galois group of  $\mathbb{F}_{q^n}$  over  $\mathbb{F}_q$  is cyclic and is generated by the Frobenius mapping  $\sigma(\alpha) = \alpha^q$ ,  $\alpha \in \mathbb{F}_{q^n}$ . A *normal* basis of  $\mathbb{F}_{q^n}$  over  $\mathbb{F}_q$  is a basis of the form  $N = \{\alpha, \alpha^q, \dots, \alpha^{q^{n-1}}\}$ , i.e. a basis that consists of the algebraic conjugates of a fixed element  $\alpha \in \mathbb{F}_{q^n}$ . Recall that an element  $\alpha \in \mathbb{F}_{q^n}$  is said to generate a normal basis over  $\mathbb{F}_q$  if its conjugates form a basis of  $\mathbb{F}_{q^n}$  as a vector space over  $\mathbb{F}_q$ . For our convenience we call a generator of a normal basis a *normal* element. A monic irreducible polynomial  $F(x) \in \mathbb{F}_q[x]$  is called *normal polynomial* or *N-polynomial* if its roots form a normal basis or, equivalently, if they are linearly independent over  $\mathbb{F}_q$ . The elements in a normal basis are exactly the roots of some *N*-polynomial. Hence, an *N*-polynomial is just another way of describing a normal basis. It is well known that such a basis always exists and any element of  $N$  is a generator of  $N$  (the normal basis theorem, see [4], Theorem 1.4.1).

The construction of *N*-polynomials over any finite field is a challenging mathematical problem. Interest in *N*-polynomials stems both from mathematical theory and practical applications such as coding theory and several cryptosystems using finite fields. The problem in general is: given an integer  $n$  and the ground field  $\mathbb{F}_q$ , construct a normal basis of  $\mathbb{F}_{q^n}$  over  $\mathbb{F}_q$ , or, equivalently, construct an *N*-polynomial in  $\mathbb{F}_q[x]$  of degree  $n$  by providing an efficient construction method.

Some results regarding constructions of special sequences  $(F_k(x))_{k \geq 0}$  of normal polynomials over  $\mathbb{F}_q$  can be found in [2, 4, 6, 7, 9] and [10, 11]. All constructions are considered as computationally easy and explicit. Cohen [3] and McNay [8] gave iterative constructions of irreducible polynomials of 2-power degree over finite

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fields of odd characteristics. Meyn [10] and Chapman [2] showed that these polynomials are  $N$ -polynomials. Another family of  $N$ -polynomials of degree  $2^k$  was suggested by Gao [4], who constructed specific sequences  $(F_k(x))_{k \geq 0}$  of  $N$ -polynomials of degree  $p^{k+2}$  over  $\mathbb{F}_p$ . In these constructions he used substitutions introduced earlier by Varshamov [13]. Kyuregyan in [6, 7] proposed a rather more general iterative technique of constructing sequences  $(F_k(x))_{k \geq 0}$  of  $N$ -polynomials of degree  $p^{k+2}$  over  $\mathbb{F}_q$  compared with the ones given by Gao [4] and Scheerhorn [11]. While in the constructions of  $N$ -polynomials over  $\mathbb{F}_{2^s}$  suggested by Gao [4] and Scheerhorn [11] the initial polynomial is a quadratic normal polynomial, in constructions suggested by Kyuregyan in [7] the initial polynomial is a normal polynomial of arbitrary degree.

In this paper, a computationally simple and explicit construction of sequences  $(F_k(x))_{k \geq 0}$  of normal polynomials and  $(F_k(x+1))_{k \geq 0}$  of self-reciprocal normal polynomials over  $\mathbb{F}_{2^s}$  is presented. For this, we will show that all members of the sequence  $(F_k(x))_{k \geq 0}$  defined by polynomials  $F_k(x) \in \mathbb{F}_{2^s}[x]$  of degree  $n2^k$  that are constructed by iterated application of the polynomial composition  $F_k(x) = x^{n2^k} F_{k-1}(\frac{x^2+x+1}{x^2})$ ,  $k \geq 0$ , for a suitable chosen initial normal polynomial  $F_0(x) \in \mathbb{F}_{2^s}[x]$  of degree  $n$ , for which the polynomial  $F_0(x+1)$  is a self-reciprocal normal polynomial, are  $N$ -polynomials and the polynomials  $F_k(x+1)$  are self-reciprocal normal polynomials over  $\mathbb{F}_{2^s}$ . Such a sequence of polynomials define a sequence of extension fields  $\mathbb{F}_{2^{sn2^k}}$  whose union is denoted by  $\mathbb{F}_{2^{sn2^\infty}} = \cup_{k \geq 0} \mathbb{F}_{2^{sn2^k}}$ .

## 2. Preliminary notes

We need the following normality results for our further study.

Let  $p$  denote the characteristic of  $\mathbb{F}_q$  and let  $n = n_1 p^e = n_1 t$ , with  $\gcd(p, n_1) = 1$ , and suppose that  $x^n - 1$  has the following factorization in  $\mathbb{F}_q[x]$ :

$$x^n - 1 = (x^{n_1} - 1)^t = (\varphi_1(x)\varphi_2(x) \cdots \varphi_r(x))^t, \quad (1)$$

where  $\varphi_i(x) \in \mathbb{F}_q[x]$  are the distinct irreducible factors of  $x^{n_1} - 1$ . For  $1 \leq i \leq r$ , let

$$\phi_i(x) = \frac{x^n - 1}{\varphi_i(x)}. \quad (2)$$

We assume that  $\phi_i(x)$  has degree  $m_i$  for  $1 \leq i \leq r$ . Furthermore, we will need Schwartz's theorem in [12] (see also [9], Theorem 4.18), which allows us to check whether an irreducible polynomial is  $N$ -polynomial.

**Proposition 2.1** ([9], Theorem 4.18) *Let  $F(x)$  be an irreducible polynomial of degree  $n$  over  $\mathbb{F}_q$  and  $\alpha$  be a root of  $F(x)$ . Let  $x^n - 1$  factor as (1) and let  $\phi_i(x)$  be as in (2). Then  $F(x)$  is  $N$ -polynomial over  $\mathbb{F}_q$  if and only if*

$$L_{\phi_i}(\alpha) \neq 0 \text{ for } i = 1, 2, \dots, r$$

where  $L_{\phi_i}(x)$  is the linearized polynomial defined by

$$L_{\phi_i}(x) = \sum_{v=0}^{m_i} t_{iv} x^{q^v} \text{ if } \phi_i(x) = \sum_{v=0}^{m_i} t_{iv} x^v.$$

A result by Jungnickel in [5] states when an element of  $\mathbb{F}_q$  is a normal bases generator. We can restate it as follows.

**Lemma 2.2** Let  $f(x) = \sum_{i=0}^n c_i x^i$  be  $N$ -polynomial of degree  $n$  over  $\mathbb{F}_q$ . Suppose  $g(x) = f(\frac{x-a}{b})$ , where  $a, b \in \mathbb{F}_q$  and  $b \neq 0$ . Then  $g(x)$  is  $N$ -polynomial if and only if  $na - b \frac{c_{n-1}}{c_n} \neq 0$ .

**Proof** Let  $n = n_1 p^e = n_1 t$ , and then by (1),  $x^n - 1$  has the following factorization in  $\mathbb{F}_q[x]$ :

$$x^n - 1 = (x^{n_1} - 1)^t = (\varphi_1(x)\varphi_2(x) \cdots \varphi_r(x))^t,$$

where  $\varphi_1(x) = x - 1$ . Set for  $i = 2, 3, \dots, r$

$$\begin{aligned} \phi_i(x) &= \frac{x^n - 1}{\varphi_i(x)} \\ &= (x - 1)^t s_i(x) \\ &= (x - 1) s'_i(x), \end{aligned}$$

where

$$s'_i(x) = (x - 1)^{t-1} s_i(x) = \sum_{v=0}^{m'_i} t'_{iv} x^v.$$

Hence,

$$\phi_i(x) = \sum_{v=0}^{m'_i} t'_{iv} x^{v+1} - \sum_{v=0}^{m'_i} t'_{iv} x^v.$$

Since  $f(x)$  is  $N$ -polynomial, then by Proposition 2.1 we have  $L_{\phi_i}(\alpha) \neq 0$  for each  $i = 1, 2, \dots, r$ , where  $\alpha$  is a root of  $f(x)$ . We need to show that  $L_{\phi_i}(a + b\alpha) \neq 0$  is also true for each  $i = 2, 3, \dots, r$ , where  $a + b\alpha$  is a root of  $g(x)$ . Since

$$L_{\phi_i}(a + b\alpha) = \sum_{v=0}^{m'_i} t'_{iv} (a + b\alpha)^{q^{v+1}} - \sum_{v=0}^{m'_i} t'_{iv} (a + b\alpha)^{q^v},$$

we have

$$\begin{aligned} L_{\phi_i}(a + b\alpha) &= a \sum_{v=0}^{m'_i} t'_{iv} + b \sum_{v=0}^{m'_i} t'_{iv} \alpha^{q^{v+1}} - a \sum_{v=0}^{m'_i} t'_{iv} - b \sum_{v=0}^{m'_i} t'_{iv} \alpha^{q^v} \\ &= b \left( \sum_{v=0}^{m'_i} t'_{iv} \alpha^{q^{v+1}} - \sum_{v=0}^{m'_i} t'_{iv} \alpha^{q^v} \right) \\ &= b L_{\phi_i}(\alpha) \neq 0, \end{aligned} \tag{3}$$

and hence, for  $g(x)$  to be an  $N$ -polynomial, it suffices to solve the condition  $L_{\phi_1}(a + b\alpha) \neq 0$ . On the other hand, we have

$$\phi_1(x) = \frac{x^n - 1}{x - 1} = x^{n-1} + x^{n-2} + \cdots + x + 1 = \sum_{i=0}^{n-1} x^i.$$

So:

$$\begin{aligned}
 L_{\phi_1}(a + b\alpha) &= \sum_{i=0}^{n-1} (a + b\alpha)^{q^i} \\
 &= \sum_{i=0}^{n-1} a + b \sum_{i=0}^{n-1} \alpha^{q^i} \\
 &= na + b \operatorname{Tr}_{q^n|q}(\alpha) = na - b \frac{c_{n-1}}{c_n},
 \end{aligned} \tag{4}$$

which is nonzero by hypothesis. This completes the proof.  $\square$

In the following propositions a family of irreducible polynomials of degree  $n2^k$  over  $\mathbb{F}_{2^s}$  is suggested. We will use them in the proof of our results.

**Proposition 2.3** ([1], Theorem 2.2) *Recalling the definitions of  $P^*$  and  $P^{*'}$ , let  $P(x) = \sum_{i=0}^n c_i x^i$  be an irreducible polynomial over  $\mathbb{F}_{2^s}$  of degree  $n$ . Then*

$$F(x) = x^{2n} P\left(\frac{x^2 + \delta_0 x + \delta_1}{x^2}\right), \quad \delta_0, \delta_1 \in \mathbb{F}_{2^s}^*$$

*is an irreducible polynomial of degree  $2n$  over  $\mathbb{F}_{2^s}$  if and only if*

$$\operatorname{Tr}_{2^s|2}\left(\frac{\delta_1}{\delta_0^2} \left(\frac{P^{*'}(0)}{P^*(0)} + n\right)\right) \neq 0.$$

**Proposition 2.4** ([1], Theorem 3.1) *Let  $P(x)$  be an irreducible polynomial of degree  $n$  over  $\mathbb{F}_{2^s}$ . Define*

$$F_0(x) = P(x),$$

$$F_k(x) = x^{n2^k} F_{k-1}\left(\frac{x^2 + x + 1}{x^2}\right) \quad k \geq 1. \tag{5}$$

*Suppose that*

$$\operatorname{Tr}_{2^s|2}\left(\frac{P'(1)}{P(1)}\right) \cdot \operatorname{Tr}_{2^s|2}\left(\frac{P^{*'}(0)}{P^*(0)} + n\right) \neq 0.$$

*Then  $(F_k(x))_{k \geq 1}$  is a sequence of irreducible polynomials over  $\mathbb{F}_{2^s}$  of degree  $n2^k$ .*

### 3. Construction of $N$ -polynomials over finite fields

In this section we establish theorems that will show how Propositions 2.3 and 2.4 can be applied to produce  $N$ -polynomials over  $\mathbb{F}_{2^s}$ .

**Theorem 3.1** *Let  $P(x) = \sum_{i=0}^n c_i x^i$ , with  $P(x) \neq x$  an  $N$ -polynomial of degree  $n$  over  $\mathbb{F}_{2^s}$  such that  $P(x+1)$  is a self-reciprocal polynomial over  $\mathbb{F}_{2^s}$ . Also let*

$$F(x) = x^{2n} P\left(\frac{x^2 + x + 1}{x^2}\right). \tag{6}$$

Then  $F(x)$  is an  $N$ -polynomial of degree  $2n$  over  $\mathbb{F}_{2^s}$ , if and only if

$$Tr_{2^s|2}\left(\frac{c_{n-1}}{c_n} + n\right) \neq 0.$$

**Proof** Recall the definition of  $Ord_{\alpha,\sigma}$ . Since  $P(x)$  is an irreducible polynomial over  $\mathbb{F}_{2^s}$ , Proposition 2.3 and the hypothesis imply that  $F(x)$  is irreducible over  $\mathbb{F}_{2^s}$ . Let  $\alpha \in \mathbb{F}_{2^{sn}}$  be a root of  $P(x)$ . Since  $P(x)$  is an  $N$ -polynomial of degree  $n$  over  $\mathbb{F}_{2^s}$  by the hypothesis,  $\alpha \in \mathbb{F}_{2^{sn}}$  is a normal element over  $\mathbb{F}_{2^s}$  and hence has order  $Ord_{\alpha,\sigma}(x) = x^n - 1$ .

Let  $n = n_1 2^e$ , where  $n_1$  is a nonnegative integer with  $gcd(n_1, 2) = 1$  and  $e \geq 0$ . For convenience we denote  $2^e$  by  $t$ . Let  $x^n - 1$  have the following factorization in  $\mathbb{F}_{2^s}[x]$ :

$$x^n - 1 = (\varphi_1(x)\varphi_2(x) \cdots \varphi_r(x))^t, \tag{7}$$

where the polynomials  $\varphi_i(x) \in \mathbb{F}_q[x]$  are the distinct irreducible factors of  $x^{n_1} - 1$ . Set

$$\phi_i(x) = \frac{(x^n - 1)}{\varphi_i(x)} = \sum_{v=0}^{m_i} t_{iv} x^v, i = 1, 2, \dots, r. \tag{8}$$

By the hypothesis  $\frac{c_{n-1}}{c_n} + n \neq 0$ , and so by Lemma 2.2,  $P(x+1)$  is a normal polynomial.

Now we proceed by proving that  $F(x)$  is a normal polynomial. Let  $\alpha_1$  be a root of  $F(x)$ .

We only need to show that the  $\sigma$ -order of  $\alpha_1$  is

$$Ord_{\alpha_1,\sigma}(x) = x^{2n} - 1.$$

Note that by (7) the polynomial  $x^{2n} - 1$  has the following factorization in  $\mathbb{F}_{2^s}[x]$ :

$$x^{2n} - 1 = (\varphi_1(x) \cdot \varphi_2(x) \cdots \varphi_r(x))^{2t},$$

where  $\varphi_i(x) \in \mathbb{F}_{2^s}[x]$  are distinct irreducible factors of  $x^{n_1} - 1$ . Let

$$H_i(x) = \frac{x^{2n} - 1}{\varphi_i(x)},$$

or

$$H_i(x) = \frac{x^{2n} - 1}{\varphi_i(x)} = (x^n + 1) \cdot \frac{x^n - 1}{\varphi_i(x)}.$$

By (8) we obtain

$$H_i(x) = (x^n + 1) \cdot \phi_i(x).$$

Hence, since  $\phi_i(x) = \sum_{v=0}^{m_i} t_{iv} x^v$ ,  $i = 1, 2, \dots, r$ , we have

$$H_i(x) = \sum_{v=0}^{m_i} t_{iv} (x^{n+v} + x^v).$$

It follows that

$$L_{H_i}(\alpha_1) = \sum_{v=0}^{m_i} t_{iv} (\alpha_1^{2^{sn}} + \alpha_1)^{2^{sv}}. \tag{9}$$

Note that, according to Proposition 2.1, to complete the proof of the theorem we only need to show that

$$L_{H_i}(\alpha_1) \neq 0 \text{ for each } i = 1, 2, \dots, r.$$

From (6), if  $\alpha_1$  is a zero of  $F(x)$ , then  $\frac{\alpha_1^2 + \alpha_1 + 1}{\alpha_1^2}$  is a zero of  $P(x)$ . It may thus be assumed that

$$\alpha = \frac{\alpha_1^2 + \alpha_1 + 1}{\alpha_1^2},$$

where  $\alpha$  is a root of  $P(x)$ . It follows that

$$\alpha + 1 = \frac{\alpha_1 + 1}{\alpha_1^2} = \frac{1}{\alpha_1} + \frac{1}{\alpha_1^2}. \quad (10)$$

Now, by (10) and observing that  $P(x)$  is an irreducible polynomial of degree  $n$  over  $\mathbb{F}_{2^s}$ , we obtain

$$\alpha + 1 = (\alpha + 1)^{2^{sn}} = \left(\frac{1}{\alpha_1} + \frac{1}{\alpha_1^2}\right)^{2^{sn}}. \quad (11)$$

It follows from (10) and (11) that

$$\left(\frac{1}{\alpha_1} + \left(\frac{1}{\alpha_1}\right)^{2^{sn}}\right)^2 = \left(\frac{1}{\alpha_1} + \left(\frac{1}{\alpha_1}\right)^{2^{sn}}\right). \quad (12)$$

It is clear that  $\left(\frac{1}{\alpha_1} + \left(\frac{1}{\alpha_1}\right)^{2^{sn}}\right) \neq 0$ .

Hence, it follows from (12) that  $\frac{1}{\alpha_1} + \left(\frac{1}{\alpha_1}\right)^{2^{sn}} = 1$ . Therefore,

$$\alpha_1^{2^{sn}} = \frac{\alpha_1}{1 + \alpha_1}. \quad (13)$$

Now by (10) and (13), we can obtain

$$\alpha_1^{2^{sn}} + \alpha_1 = \frac{1}{\alpha + 1}. \quad (14)$$

Thus, by (9) and (14), we have

$$L_{H_i}(\alpha_1) = \sum_{v=0}^{m_i} t_{iv} \left(\frac{1}{\alpha + 1}\right)^{2^{sv}}. \quad (15)$$

Since  $\frac{1}{\alpha + 1}$  is a zero of the normal polynomial  $(P(x + 1))^*$ , therefore  $L_{H_i}(\alpha_1) \neq 0$ . Hence,  $F(x)$  is a normal polynomial of degree  $2n$  over  $\mathbb{F}_{2^s}$ , and the proof is completed.  $\square$

#### 4. Recurrent methods for constructing normal polynomials

In this section we describe a computationally simple and explicit recurrent method for constructing higher degree normal polynomials over finite fields  $\mathbb{F}_{2^s}$  starting from a normal polynomial. We begin by establishing the following theorem.

**Theorem 4.1** Let  $P(x) = \sum_{i=0}^n c_i x^i$ , with  $P(x) \neq x$  an  $N$ -polynomial of degree  $n$  over  $\mathbb{F}_{2^s}$  such that  $P(x+1)$  is a self-reciprocal polynomial over  $\mathbb{F}_{2^s}$ . Define

$$F_0(x) = P(x),$$

$$F_k(x) = x^{n2^k} F_{k-1}\left(\frac{x^2 + x + 1}{x^2}\right) \quad k \geq 1. \tag{16}$$

Then  $(F_k(x))_{k \geq 0}$  and  $(F_k(x+1))_{k \geq 0}$  are the sequences of  $N$ -polynomials and self-reciprocal  $N$ -polynomials of degree  $n2^k$  over  $\mathbb{F}_{2^s}$ , respectively, if and only if

$$\text{Tr}_{2^s|2}\left(\frac{P'(1)}{P(1)}\right) \cdot \text{Tr}_{2^s|2}\left(\frac{c_{n-1}}{c_n} + n\right) \neq 0,$$

where  $P'(1)$  is the formal derivative of  $P(x)$  at point 1.

**Proof** It is easy to check that the polynomial  $F_k(x+1)$ , for each  $k \geq 1$ , is self-reciprocal by using the definitions. According to Proposition 2.4 for each  $k \geq 1$ ,  $F_k(x)$  is an irreducible polynomial over  $\mathbb{F}_{2^s}$ . Consequently,  $(F_k(x+1))_{k \geq 0}$  is a sequence of irreducible polynomials over  $\mathbb{F}_{2^s}$ . The proof of normality of the irreducible polynomial  $F_k(x)$  for each  $k \geq 1$  is done by mathematical induction on  $k$ .

For  $k = 1$ ,  $F_1(x)$  is a normal polynomial according to Theorem 3.1.

For  $k \geq 2$ , we show that  $F_k(x)$  is also a normal polynomial. To this end we need to show that the hypothesis of Theorem 3.1 is satisfied. However, by induction hypothesis, we have  $F_{k-1}(x)$  as a normal polynomial and  $F_{k-1}(x+1)$  as a self-reciprocal polynomial. Thus, by Theorem 3.1,  $F_k(x)$  is a normal polynomial if and only if

$$\text{Tr}_{2^s|2}\left(\frac{F_{k-1}^{*'}(0)}{F_{k-1}^*(0)} + 2^{k-1}n\right) \neq 0,$$

or

$$\text{Tr}_{2^s|2}\left(\frac{F_{k-1}^{*'}(0)}{F_{k-1}^*(0)}\right) \neq 0.$$

However, from (16), we have

$$\begin{aligned} F_{k-1}^*(x) &= x^{n2^{(k-1)}} F_{k-1}\left(\frac{1}{x}\right) \\ &= x^{n2^{(k-1)}} \left(\frac{1}{x}\right)^{n2^{(k-1)}} F_{k-2}\left(\frac{\left(\frac{1}{x}\right)^2 + \left(\frac{1}{x}\right) + 1}{\left(\frac{1}{x}\right)^2}\right) \\ &= F_{k-2}(x^2 + x + 1). \end{aligned} \tag{17}$$

So

$$F_{k-1}^*(0) = F_{k-2}(1) \tag{18}$$

and

$$F_{k-1}^{*'}(0) = F'_{k-2}(1). \tag{19}$$



On the other hand:

$$F'_{k-1}(x) = x^{n2^{(k-1)}-2} F'_{k-2}\left(\frac{x^2+x+1}{x^2}\right). \quad (20)$$

So

$$F'_{k-1}(1) = F'_{k-2}(1). \quad (21)$$

Using (19) and (21), we get

$$F_{k-1}^{*'}(0) = P'(1). \quad (22)$$

Obviously by (16)

$$F_{k-1}(1) = F_{k-2}(1). \quad (23)$$

So (18) and (23) imply that

$$F_{k-1}^*(0) = P(1). \quad (24)$$

Hence, by (22) and (24) we obtain

$$Tr_{2^s|2}\left(\frac{F_{k-1}^{*'}(0)}{F_{k-1}^*(0)}\right) = Tr_{2^s|2}\left(\frac{P'(1)}{P(1)}\right), \quad (25)$$

which is not equal to zero by the hypothesis of the theorem and so  $(F_k(x))_{k \geq 0}$  is a sequence of  $N$ -polynomials of degree  $n2^k$  over  $\mathbb{F}_{2^s}$ . Finally, we note that by Lemma 2.2, for every  $k \geq 1$ ,  $F_k(x+1)$  is an  $N$ -polynomial if and only if  $F_k^{*'}(0) \neq 0$ . Thus, (22) and the hypothesis of the theorem imply that  $(F_k(x+1))_{k \geq 0}$  is a sequence of self-reciprocal  $N$ -polynomials of degree  $n2^k$  over  $\mathbb{F}_{2^s}$ . The theorem is proved.  $\square$

**Example 4.2** Consider the normal polynomial  $P(x) = x^2 + x + 1$  over  $\mathbb{F}_2$ . It is easy to see that the assumptions of Theorem 4.1 are fulfilled. Therefore, the composite polynomials

$$F_1(x) = x^4 P\left(\frac{x^2+x+1}{x^2}\right) = x^4 + x^3 + 1$$

and

$$\begin{aligned} F_2(x) &= x^8 F_1\left(\frac{x^2+x+1}{x^2}\right) \\ &= x^8 + x^7 + x^5 + x^4 + x^3 + x^2 + 1 \end{aligned}$$

are normal polynomials over  $\mathbb{F}_2$ . Furthermore, the polynomials

$$F_1(x+1) = x^4 + x^3 + x^2 + x + 1$$

and

$$F_2(x+1) = x^8 + x^7 + x^6 + x^4 + x^2 + x + 1$$

are self-reciprocal normal polynomials over  $\mathbb{F}_2$ . Obviously, Theorem 4.1 describes a computationally simple and explicit recurrent method for constructing normal and self-reciprocal normal polynomials, so computing the normal and self-reciprocal normal polynomials  $F_k(x)$  and  $F_k(x+1)$ , respectively, for  $k \geq 3$  is not a complex problem.

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